

Efficient Auditing Mechanism for Data Storage in Cloud Using Abs

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Abstract: Cloud computing's multi-tenancy feature, which provides privacy, security and access control challenges, because of sharing of physical resources among untrusted tenants. In order to achieve safe storage, policy based file access control, policy based file assured deletion and policy based renewal of a file stored in a cloud environment. A suitable encryption technique with key management should be applied before outsourcing the data. Existing system attribute based encryption is a new cryptographic primitive which provides a promising tool for addressing the problem of secure and fine-grained data sharing and decentralized access control. Key policy attribute based encryption (KP-ABE) is an important type of ABE, which enables senders to encrypt messages under a set of attributes and private keys are associated with access structures that specify which ciphertexts the key holder will be allowed to decrypt. In this paper (ABS) attribute based signature new algorithm can be easily implemented and greatly raise the computational efficiency of original ABS algorithm. A verifier will be convinced of the fact that whether the signer's attributes satisfy the predicate while remaining completely ignorant of the identity of the signer. ABS is much useful in a wide range of applications including private access control for ad hoc networks, attribute based messaging, etc.

Keywords: Attribute based signature, Cloud storage, Cloud service provider, Data owner, Encryption.

I. INTRODUCTION

Nowadays cloud computing is a rationally developed technology to store data from more than one client. Cloud computing is an environment that enables users to remotely store their data. Remote backup system is the advanced concept which reduces the cost for implementing more memory in an organization. It helps enterprises and government agencies reduce their financial overhead of data management. They can archive their data backup remotely to third party cloud storage providers rather than maintain data center on their own. An individual; or an organization may not require purchasing the needed storage device. Instead they can store their data backups to the cloud and archive their data to avoid any information loss in case of hardware/software failures. Much of the data stored in clouds is highly sensitive, for example/medical records and social networks. Even cloud storage is more flexible, how the security and privacy are available for the outsourced data becomes a serious concern. In one hand, the user should authenticate it before initiating any transaction, and on the other hand, it must be ensured that the cloud does not tamper with the data that is outsourced. User privacy is also required so that the cloud or other users do not know the identity of the user. The cloud can hold the user accountable for the data it outsources, and likewise, the cloud is itself accountable for the service it provides. The validity of the user who stores the data is also verified. Apart from the technical solutions to ensure security and privacy, there is also a need for law enforcement. Access control in clouds is gaining attention because it is important that only authorized users have access to valid services. A huge amount of information is being stored in the cloud, and much of this sensitive.

Care should be taken to ensure access control of this sensitive information which can often be related to health, important documents (as in Google Docs or Drop Box) or even personal information (as in social networking). It is just not enough to store the contents securely in the cloud but it might also be necessary to ensure anonymity of the user. For example, a user would like to store some sensitive information but does not want to be recognized. Thus a user might want to post a comment on an article. But does not want his/her identity to disclose. However, the user should be able to prove to the other user that he/she is a valid user who stored the information without revealing the identity. There are cryptographic protocols like ring signatures [13], mesh signatures [14], group signature [15]. This can be used in these situations. Ring signature is not a feasible option for clouds where there are a large number of users. Group signature assumes the preexistence of a group which might not be possible in clouds. After comparing the drawbacks of all the cryptographic protocols mentioned above, a new protocol known as attribute based signature (ABS) has been proposed in this paper. ABS was proposed to help identify the users as an authorized one, without revealing its identity. Other users or the cloud can verify the user and the validity of the message stored. ABS can be combined with ABE to authenticated access control without disclosing the identity of the user to the cloud.

II. LITERATURE SURVEY

Attribute based signature were first introduced by Magji, Prabhakaran, and Rosales (2008) as a way to let a signature attest not to be the identity of the individual who

endorsed a message, but instead to a (possibly complex) claim regarding the attribute she processes. They constructed an ABS scheme that supports a powerful set of predicates, namely, any predicate consists of AND, OR, and threshold gates. However, the security of their scheme is weak as their construction is only proved in the generic group model. Since then, there have been lots of works on this subject (Escala; Herants; Morillo, 2011; Khader, 2007a; 2007b; Au; Susilo; Xie; Ren, 2010; Li; Kim, 2007; 2010; Maji; Prabharan; Rosulek, 2011; Okamoto; Takashima, 2011; Shahandashti; Safavinaili, 2009). Recently, Maji, Prabharan, and Rosulek (2011) presented an ABS scheme which is proven secure in the standard model, but it is much less efficient and more complicated than the scheme in (Maji; Prabharan; Rosulek, 2008), since it employs the Groth-Sahai NIZK protocol (2008) as building blocks. Okamoto and Takashima (Okamoto; Takashima, 2011) presented a fully secure attribute-based signature (ABS) scheme in the standard model. The admissible predicates of the scheme support non-monotone predicates. Escala, Herranz, and Morillo (2011) proposed a fully secure attribute-based signature (ABS) scheme in the standard model. This scheme supports an additional property of revocability, so that an external judge can break the anonymity of a signature when necessary. Another related notion to ABS is fuzzy identity-based signature, which was proposed and formalized in (Shanqing; Yingpei, 2008; Tang; Cao; Dong, 2011). It allows a user with identity ω to issue a signature which could be verified with identity ω' if and only if ω and ω' are within distance judged by some metric. However, this kind of signature does not consider the anonymity for the signer.

Our contribution

Maji, Prabharan, Rosulek (Maji; Prabharan; Rosulek, 2008; 2011) and Okamoto, Takashima (2011) pointed that the future work of ABS, on the theoretical front, is to base the security of ABS on a standard hardness assumption, while still reserve the efficiency for the most part. In this paper, we attempt to propose such an ABS scheme which is secure in the standard model based on the decisional parallel bilinear Diffie-Hellman exponent assumption (Water, 2011) which is more practical than the generic group model of (Maji; Prabharan; Rosulek, 2008).

The proposed ABS scheme is efficient and practical, we compare our scheme with the existing ABS scheme in the standard model. Maji, Prabharan, and Rosulek's (2011) (two typical instantiations), Okamoto and Takashima's (2011), and Escala, Herranz, and Morillo's (2011), as well as the ABS scheme in the generic group model (Maji; Prabharan; Rosulek, 2008) (as a benchmark). All of these schemes can be implemented over a pairing group and the size of group element is about the size of zp (e.g. 256 bits). Our construction is inspired by the attribute-based encryption scheme (ABE) of Water (2011). Roughly speaking, a secret signing key sks with attribute set S corresponds to a secret decryption key sks in ABS (Water, 2011). No counterpart of signature σ in our construction exists in the ABE (Water, 2011).

III. PROPOSED WORK

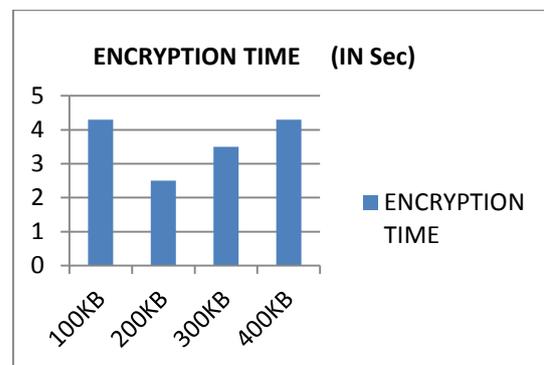
Cloud computing, as an emerging computing paradigm, enables users to remotely store their data into a cloud so as to enjoy scalable services on-demand. Especially for small and medium-sized enterprises with limited budgets, they can collaborate, and the like. However, allowing cloud services providers (CSOs), which are not in the same trusted domains as enterprise users to make use of confidential data, may raise potential security and privacy issues. To make the sensitive user data confidential against untrusted CSPs, a natural way is to apply cryptographic approach. By disclosing decryption keys only to authorized users. However, when enterprise users outsource confidential data for sharing on cloud servers, the adopted encryption system should not only support fine-grained access control, but also provide high performance, full delegation, and scalability, so as to best serve the needs of accessing data anytime and anywhere, delegating within enterprises. And achieving a dynamic set of users.

IV. EXPERIMENTAL RESULT

The hardware used to carry out this experiment is an Intel Core i3 processor with 2GB RAM. The program is written in C# programming; it is used for compilation and execution purposes.

Encryption

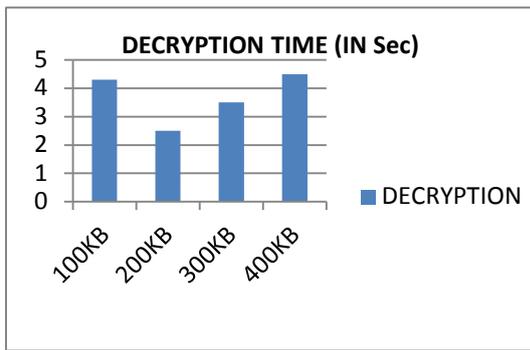
In the first graph, we show the comparison of time required for encryption using the ABS algorithm.



This graph gathers FIVE different files with different sizes to measure the performance metrics using encryption time. The time will vary with different file size and type of the file.

Decryption

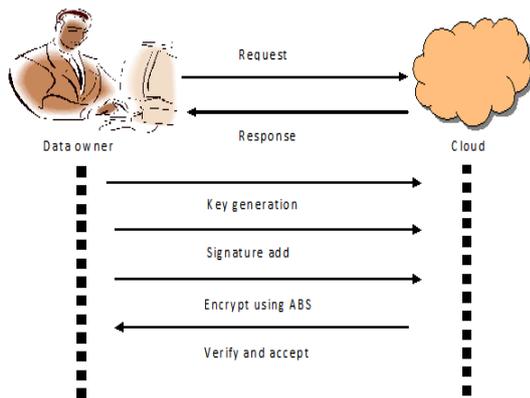
In the second graph, we show the comparison of time for the ABS algorithm. This graph gathers FIVE different encrypted files received from the sender and decrypted using the ABS public key. In this process, we measure the performance metrics using decryption time. The time varies with different encrypted file formats.



Summary

So the detailed analysis between the encryption and decryption time with different file formats. after analyzing the results, that shows the ABS algorithm time dependant on file format, also this work conclude ABS is efficiency handle encryption and decryption with performance matrices. In ABS the decryption time is much less from the taken in encryption method.

METHODOLOGY USED



Cloud Computing

Cloud computing is becoming a well-known buzzword nowadays. as a brand new infrastructure to offer services, cloud computing system have many superiorities in comparing to those existed traditional service provisions, such as reduced upfront investment, expected performance, high availability, infinite scalability, tremendous fault-tolerance capability and so on and consequently chased by most of the it companies, such as google, Amazon, Microsoft, salesforce.com. Based on their overwhelming predominance in traditional service provisions and capital accumulation, most of these it companies have more chance to adapt their service into such a new environment earlier, say cloud computing system. on the other hand ,a large number of new companies are spawned with competitive services relayed on those provided cloud computing system. in terms of their provisions ,we divide those service into servel categories are software as service(SaaS),platform as a service (PaaS),identity and policy management as a service(ipmaas),network as a service (Naas).infrastructure as a service(IaaS).detailed analysis to these services are provide the corresponding service categories.

Cloud auditing

A party that can conduct independent assessment of cloud services, information system operation. Performance and security of the cloud implementation. a cloud auditor can evaluate the services provided by a cloud provider in terms of security controls, privacy impact, performance, etc. for security auditing a cloud auditor can make an assessment of the security controls in the information system to determine the extent to which the controls are implemented correctly. Operating as intended, and producing the desired outcome with respect to meeting the security requirements for the system. public auditing is the service which is used to ensure integrity of the data stored on the cloud user's computation resources. To perform the auditing task the TPA known as third party audit the stored data on cloud.TPA verify the correctness of the cloud data on demand without retrieving a copy of the whole data. The TPA, has expertise and capabilities that can periodically check the integrity of all the data stored which provides a much more easier and affordable way for the users to ensure their storage correctness in the cloud.

ABS: attribute based signature

attribute based signature (ABS) is a new and versatile cryptographic primitive proposed in [1,2].in which a signature attests not to the identity of the individual who endorsed a messgae, but instead to a (possibly complex) claim regarding the attribute authority.ABS offer an unforgeability guarantee for the verifier, that the signature was produced by a single party whose attribute satisfy the policy being made;ie not by collusion of individuals who pooled their attributes together. Although the scheme in [2] reveals the set of attributes satisfying the policy, subsequent research of ABS offer an attribute signer privacy guarantee for the signer. That is ,a legitmate signer remains anonymous without the fear of revocation and is indistinguishable among all the users whose attribute satisfying the policy specified in the signature.ABS is useful in many important applications such as anonymous authentication and attribute based messaging system. In ABS in which a signer is defined by a set of attribute instead of a single string reprinting the signer's identity. In ABS, a user obtains a set of attributes from one pr multiple attribute authorities. An attribute based signature assures the verifier that a signer, whole set of attribute satisfies a (possibly) complex predicate, has endorsed the message. The following example illustrates the concept. Suppose we have the following predicate: professor **OR** (((Biology Department **OR** Female) **OR** above 50 years old) **AND** university A) Alice's attributes are (University A, female).Bob's attribute are (above 50 years old, professor).although their attribute are quite are quite different. it is clear that alice and Bob can generate a signature on this predicate, and such a signature releases no information regarding the attribute or identity of the signer,i.e.Alice or Bob, except that the attribute of the signer satisfile the predicate.

ALGORITHM STEPS

System initialization

Select a prime q and groups g_1, g_2 . define the mapping $e: G_1 \times G_2 \rightarrow G$. let $H \rightarrow$ hash function. Let $A_0 = h_0$ and (T_{sig}, T_{ver}) mean T_{sig} is the private key with which a message is signed and T_{Ver} is the public key used for verification.

1. Secret for the trustee is $TSK = (a, T_{Sig})$
2. Public key is $TPK = (G_1, G_2, H, g_1, A_0, h_0, h_1, \dots, h_{tmax}, g_2, T_{Ver})$

User registration

for a user with identity U_u the KDC draw at random $K_{base} \in G$. $LETKO = K_{BASE}^{1/a_0}$. Token is generated

$$as: \gamma = (u, k_{base}, K_o, \rho)$$

kdc setup

Choose $a, b \in \mathbb{Z}_q^*$ randomly and compute

$$A_{ij} = g_j^a \text{ and } B_{ij} = h_j^b$$

1. Private key of KDC is $ASK[i] = (a, b)$
2. Public key of KDC is $APK [I] = (A_{ij}, B_{ij} |_{j \in [1, tmax]})$

Attribute generation

The token verification algorithm verifier the signature contained in γ using the signature verification key T_{Ver} in

TPK . this algorithm extract K_{base} from γ and $\Lambda K_x =$

$$K_{base}^{1/(a+bx)}$$

the key K_x can be checked for consistency using algorithm

$ABS.keycheck(TP, APK[I], \gamma, K_x)$ Which checks

$$e^{\Lambda}(K_x, A_{ij} B_{ij}^x) e^{\Lambda}(K_{base}, h_j)$$

Sign

The algorithm ABS

$.sign(TPK, \{APK[i]: i \in AT[U]\}, \gamma, \{K_x: X \in J_u\}, MSG)$ has input the public key of the trustee, the secret key for the message to be signed and the policy claim y .

1. Compute $\mu = H(MSG || y)$
2. Choose r_o and r_i then compare Y, W
3. $Y = K_{base}^{r_o}, S_i = (k_i^{vi})^{r_o} \cdot (g_2 g_1^{\mu})^{r_i} (\forall i \in ju)$
4. $W = K_o^{r_o}, p_j = \prod_{I \in AT[U]} (A_{ij} B_{ij}^{\pi' i})^{M_{ij} r_i} (\forall j \in [t])$
5. Signature: $\sigma = Y, W, S_1, S_2, \dots, S_t, p_1, \dots, p_t$

Verify

This uses the algorithm ABS. $verify(TPK, \sigma = (Y, W, S_1, S_2, \dots, S_t, p_1, p_2, \dots, p_t) MSG, Y)$

1. Compute $\mu = H(MSG || y)$
2. if $Y=1 \rightarrow AB.verify = 0 \rightarrow false$
3. or else check $e^{\Lambda}(W, A_o) = e^{\Lambda}(Y, h_o)$

V. CONCLUSION & FUTURE WORK

we introduced a new encryption primitive called attribute based signature (ABS), and presented a construction that provably meets that requirements of ABS. the security proof we provide is in the key generation and encrypt using ABS. the construction is fairly efficient. for reasonably complex claim predicates. ABS scheme supports multiple attribute authorities and multiple signature-trustees, who need not trust each other. IN THIS future work, on the theoretical front, one would like to base the security on a standard hardness assumption, without sacrificing too much of the efficiency. the evident drawback of relying on security in the generic group model is in evaluating the underlying group used by a candidate implementation. in fact, even the "standard" hardness assumption, as the algorithmic study of such groups is relatively in its infancy.

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